# ANALYSIS AND OPTIMIZATION OF MARKOV HM-NETWORKS WITH STOCHASTIC INCOMES FROM TRANSITIONS BETWEEN THEIR STATES 

Michal Matalycki ${ }^{1}$, Katarzyna Koluzaewa ${ }^{2}$<br>${ }^{1}$ Institute of Mathematics, Czestochowa University of Technology, Poland<br>${ }^{2}$ Faculty of Mathematics and Computer Science, Grodno State University, Bielarus<br>koluzaeva@gmail.com


#### Abstract

Expressions for expected incomes and variations of incomes in systems of Markov HM-networks, when service rates of messages are depending or not depending of network's states are obtained. The case when incomes of transitions between network's states are random variables with the set moments of first two orders is considered. Solution of some optimization problems for HM-networks are resulted.


## Introduction

It is known that functioning of any Markov queueing network $(\mathrm{QN})$ is described with a help of Markovian chain with continuous time. In Howard's work [1] the conception of Markovian chains with incomes that were constants was introduced and there were proposed to use method of Laplace transformation and method of z-transformation for analysis of such chains with small number of states. This concept has laid down in a basis of definition of Markov QN with incomes that were examined in works [2-4] at first. Later open and closed networks with central queueing system (QS) were investigated in cases when a) incomes from transitions between network's states depend on states and time [5-7] or b) incomes are random variables (RV) with the set moments of the first and the second orders [8, 9]. Recently QN with incomes refer to as HM (Howard-Matalytski)-networks. In the present article expressions in case b) were obtained for expected incomes and income's variations in systems of open exponential HM-networks and Jackson HMnetworks of arbitrary architecture.

Let us examine open exponential QN with one type messages that consist of $n$ QS $S_{1}, S_{2}, \ldots, S_{n}$. Vector $k(t)=(k, t)=\left(k_{1}, k_{2}, \ldots, k_{n}, t\right)$ is state of network where $k_{i}$ is number of messages in system $S_{i}$ at the moment $t, t \in\left[t_{0},+\infty\right), i=\overline{1, n}$. For unification of designation let us introduce system $S_{0}$ (outside medium) from which the Poisson flow of messages with arrival rate $\lambda$ comes into network. At first we will
examine the case when service rate of messages $\mu_{i}$ in system $S_{i}$ doesn't depend on number of messages in it, $i=\overline{1, n}$. Let $p_{0 j}$ is probability of message entry from system $S_{0}$ to system $S_{j}, \sum_{j=1}^{n} p_{0 j}=1 ; p_{i j}$ is probability of message transition after its service in system $S_{i}$ to system $S_{j}, \sum_{j=0}^{n} p_{i j}=1, i=\overline{1, n}$. Message during its transition from system $S_{i}$ to system $S_{j}$ brings to system $S_{j}$ some random income and income of system $S_{i}$ descend on this value correspondently, $i, j=\overline{0, n}$. It is needed to find expected (average) incomes of network's systems during time $t$ considering that network's state in the starting time $t_{0}$ is known. At first we will suppose that network is functioning in the term of high capacity, i.e. $k_{i}(t)>0, \forall t \in\left[t_{0},+\infty\right)$, $i=\overline{1, n}$.

## 1. Expected incomes of network systems

Let us examine dymamics of income's change of some network's system $S_{i}$. Let us denote its income at the moment $t$ as $V_{i}(t)$. The income of this system at some moment $t_{0}$ is equal to $v_{i 0}$. We will be interested in system's income $V_{i}\left(t_{0}+t\right)$ at the moment $t_{0}+t$. To a finding of the income we shall apply a following technique. Let us divide interval $\left[t_{0}, t_{0}+t\right]$ into $m$ equal portions with length $\Delta t=\frac{t}{m}$, and consider that $m$ is big. For finding of system's income Let us write probability of all events that can occur at the $l$-th time interval, $l=\overline{1, m}$. The next situations are possible: a) message from outside will arrive to system $S_{i}$ with probability $\lambda p_{0 i} \Delta t+o(\Delta t)$ and will bring income of size $r_{0 i}$ to it, where $r_{0 i}$ - RV with distribution function (d.f.) $F_{0 i}(x) ;$ b) message from system $S_{i}$ will transmit with probability $\mu_{i} p_{i 0} \Delta t+o(\Delta t)$ into outside medium and system's income will decrease on value $R_{i 0}$, where $R_{i 0}$ - RV with d.f. $F_{i 0}(x)$; c) message from system $S_{j}$ will transmit with probability $\mu_{j} p_{j i} \Delta t+o(\Delta t)$ into system $S_{i}$ and income of system $S_{i}$ will increase on value $r_{j i}$ and income of system $S_{j}$ will decrease on this value, $j=\overline{1, n}, j \neq i$, where $r_{j i}-\mathrm{RV}$ with d.f. $F_{1 j i}(x)$; d) message from system $S_{i}$ will transmit with probability $\mu_{i} p_{i j} \Delta t+o(\Delta t)$ into system $S_{j}$ and income of system $S_{j}$
will increase on value $R_{i j}$ and income of system $S_{i}$ will decrease on this, where $R_{i j}$ - RV with d.f. $F_{2 i j}(x), i, j=\overline{1, n}$; e) change of system $S_{i}$ state won't occur with probability $1-\left(\lambda p_{0 i}+\mu_{i}\right) \Delta t+o(\Delta t)$ at the time interval $\Delta t$.
It's evident that $r_{j i}=R_{j i}$ with probability $1, i, j=\overline{1, n}$, i.e.

$$
\begin{equation*}
F_{1 i j}(x)=F_{2 i j}(x), i, j=\overline{1, n} \tag{1}
\end{equation*}
$$

Besides system $S_{i}$ increases its income on value $r_{i} \Delta t$ during every small time interval $\Delta t$ due to interest on money that are in it; let us suppose that $r_{i}$ is RV with d.f. $F_{i}(x), i=\overline{1, n}$. We will suppose also that $\mathrm{RV} r_{j i}, R_{i j}, r_{0 i}, R_{i 0}, r_{i}$ are independent in pairs, $i, j=\overline{1, n}$.

Let $d_{i l}(\Delta t)$ is income change of $i$-th QS at the $l$-th time interval of length $\Delta t$. Then from all that was saying before it follows

$$
d_{i l}(\Delta t)=\left\{\begin{array}{l}
r_{0 i}+r_{i} \Delta t \quad \text { with probability } \lambda p_{0 i} \Delta t+o(\Delta t),  \tag{2}\\
-R_{i 0}+r_{i} \Delta t \quad \text { with probability } \mu_{i} p_{i 0} \Delta t+o(\Delta t), \\
r_{j i}+r_{i} \Delta t \quad \text { with probability } \mu_{j} p_{j i} \Delta t+o(\Delta t), j=\overline{1, n}, j \neq i, \\
-R_{i j}+r_{i} \Delta t \quad \text { with probability } \mu_{i} p_{i j} \Delta t+o(\Delta t), j=\overline{1, n}, j \neq i, \\
r_{i} \Delta t \quad \text { with probability } 1-\left(\lambda p_{0 i}+\mu_{i}\right) \Delta t+o(\Delta t) .
\end{array}\right.
$$

Income of system $S_{i}$ equals

$$
\begin{equation*}
V_{i}(t)=v_{i 0}+\sum_{l=1}^{m} d_{i l}(\Delta t) \tag{3}
\end{equation*}
$$

Let us introduce definition for expectation values (e.v.) correspondently:

$$
\begin{gather*}
M\left\{r_{j i}\right\}=\int_{0}^{\infty} x d F_{1 j i}(x)=a_{j i}, M\left\{R_{i j}\right\}=\int_{0}^{\infty} x d F_{2 i j}(x)=b_{i j}, M\left\{r_{i}\right\}=\int_{0}^{\infty} x d F_{i}(x)=c_{i} \\
M\left\{r_{0 i}\right\}=\int_{0}^{\infty} x d F_{0 i}(x)=a_{0 i}, M\left\{R_{i 0}\right\}=\int_{0}^{\infty} x d F_{i 0}(x)=b_{i 0}, i, j=\overline{1, n} \tag{4}
\end{gather*}
$$

in view of (1)

$$
\begin{equation*}
a_{j i}=b_{j i}, i, j=\overline{1, n} \tag{5}
\end{equation*}
$$

Let us derive expression for e.v. of income of system $S_{i}$ at the moment $t$. According to formula for e.v. of discrete RQ

$$
\begin{equation*}
M\left\{d_{i l}(\Delta t)\right\}=\left(\lambda a_{0 i} p_{0 i}+\sum_{j=1}^{n} \mu_{j} a_{j i} p_{j i}-\mu_{i} \sum_{j=0}^{n} b_{i j} p_{i j}+c_{i}\right) \Delta t+o(\Delta t), i=\overline{1, n} \tag{6}
\end{equation*}
$$

As it follows from (3) taking into account $m \Delta t=t$, when $\Delta t \rightarrow 0$

$$
\begin{equation*}
v_{i}(t)=M\left\{V_{i}(t)\right\}=v_{i 0}+\left[c_{i}+\lambda a_{0 i} p_{0 i}+\sum_{j=1}^{n} \mu_{j} a_{j i} p_{j i}-\mu_{i} \sum_{j=0}^{n} b_{i j} p_{i j}\right] t \tag{7}
\end{equation*}
$$

For expected income for whole network considering (5), we have:

$$
\begin{equation*}
M\{W(t)\}=\sum_{i=1}^{n} v_{i}(t)=\sum_{i=1}^{n}\left[v_{i 0}+\left(c_{i}+\lambda a_{0 i} p_{0 i}-\mu_{i} b_{i 0} p_{i 0}\right) t\right] \tag{8}
\end{equation*}
$$

Let us notice that common expected network's income doesn't depend on $r_{i j}, R_{i j}, i, j=\overline{1, n}$, as that incomes absorb each other (if message transmits from one network's system to another then income of the last one increases on some value and income of the first QS decreases on the same value).

## 2. Variations of network systems incomes

From expressions (2), (3) it follows that

$$
M^{2}\left\{V_{i}(t)-v_{i 0}\right\}=\left[c_{i}+\lambda a_{0 i} p_{0 i}-\mu_{i} b_{i 0} p_{i 0}+\sum_{j=1}^{n}\left(\mu_{j} a_{j i} p_{j i}-\mu_{i} b_{i j} p_{i j}\right)\right]^{2} t^{2}
$$

Let us introduce denotations:

$$
\begin{equation*}
M\left\{{r_{0 i}}^{2}\right\}=a_{20 i}, \quad M\left\{R_{i 0}^{2}\right\}=b_{2 i 0}, M\left\{r_{j i}^{2}\right\}=a_{2 j i}, \quad M\left\{R_{i j}^{2}\right\}=b_{2 i j}, i, j=\overline{1, n} \tag{9}
\end{equation*}
$$

Let us consider expression

$$
\begin{align*}
M\left(V_{i}(t)\right. & \left.-v_{i 0}\right)^{2}=M\left(v_{i 0}+\sum_{l=1}^{m} d_{i l}(\Delta t)-v_{i 0}\right)^{2}=M\left(\sum_{l=1}^{m} d_{i l}(\Delta t)\right)^{2}= \\
& =\sum_{l=1}^{m} M d_{i l}^{2}(\Delta t)+\sum_{\substack{l=1}}^{m} \sum_{\substack{j=1 \\
j \neq l}}^{m} M\left(d_{i l}(\Delta t) d_{i j}(\Delta t)\right), i=\overline{1, n} \tag{10}
\end{align*}
$$

Taking into account (2)-(9), we have:

$$
\begin{gather*}
M\left\{d_{i l}^{2}(\Delta t)\right\}=M\left[\left(r_{0 i}^{2}+2 r_{0 i} r_{i} \Delta t+r_{i}^{2}(\Delta t)^{2}\right) \lambda p_{0 i} \Delta t\right]+ \\
+M\left[\left(R_{i 0}^{2}-2 R_{i 0} r_{0} \Delta t+r_{i}^{2}(\Delta t)^{2}\right) \mu_{i} p_{i j} \Delta t\right]+\sum_{\substack{j=1 \\
j \neq i}}^{n} M\left[\left(r_{j i}^{2}+2 r_{j i} r_{i} \Delta t+\right.\right. \\
\left.\left.+r_{i}^{2}(\Delta t)^{2}\right) \mu_{j} p_{j i} \Delta t\right]+\sum_{\substack{j=1 \\
j \neq i}}^{n} M\left[\left(R_{i j}^{2}-2 R_{i j} r_{i} \Delta t+r_{i}^{2}(\Delta t)^{2}\right) \mu_{i} p_{i j} \Delta t\right]+ \\
+M r_{i}^{2}(\Delta t)^{2}\left[1-\left(\lambda p_{0 i}+\mu_{i}\right) \Delta t\right]+o(\Delta t)= \\
=\left[\lambda a_{20 i} p_{0 i}+\mu_{i} b_{2 i 0} p_{i 0}+\sum_{\substack{j=1 \\
j \neq i}}^{n}\left(\mu_{j} a_{2 j i} p_{j i}+\mu_{i} b_{2 i j} p_{i j}\right)\right] \Delta t+o(\Delta t), i=\overline{1, n} \tag{11}
\end{gather*}
$$

Besides from independence of $d_{i l}(\Delta t)$ and $d_{i j}(\Delta t), j \neq i$, considering (2) it follows that $M\left(d_{i l}(\Delta t) d_{i q}(\Delta t)\right)=o(\Delta t)^{2}$. So as it's fallowed from (10), (11) and $m \Delta t=t$, for variation of income of $i$-th QS when $\Delta t \rightarrow 0$ we obtain the following expression

$$
\begin{align*}
& \left.D V_{i}(t)=D\left(V_{i}(t)-v_{i 0}\right)=M^{\{ }\left(V_{i}(t)-v_{i 0}\right)^{2}\right\}-M^{2}\left\{\left(V_{i}(t)-v_{i 0}\right)\right\}= \\
& =\left[\lambda a_{20 i} p_{0 i}+\mu_{i} b_{2 i 0} p_{i 0}+\sum_{\substack{j=1 \\
j \neq i}}^{n}\left(\mu_{j} a_{2 j i} p_{j i}+\mu_{i} b_{2 i j} p_{i j}\right)\right] t+ \\
& +\left[\mu_{i} b_{i 0} p_{i 0}-\lambda a_{0 i} p_{0 i}-c_{i}+\sum_{j=1}^{n}\left(\mu_{i} a_{i j} p_{i j}-\mu_{j} a_{j i} p_{j i}\right)\right] t^{2}, i=\overline{1, n} \tag{12}
\end{align*}
$$

## 3. Expected incomes of Jackson HM-network systems

Let us consider now the case when service rate of messages $\mu_{i}\left(k_{i}\right)$ in system $S_{i}$ depends on number of messages that are present in it, $i=\overline{1, n}$. So we take off limitations that network must functioning in the term of high capacity. In our case expression (2) will look as

$$
d_{i l}(\Delta t)=\left\{\begin{array}{c}
r_{0 i}+r_{i} \Delta t \quad \text { with probability } \lambda p_{0 i} \Delta t+o(\Delta t) \\
-R_{i 0}+r_{i} \Delta t \quad \text { with probability } \mu_{i}\left(k_{i}(l)\right) u\left(k_{i}(l)\right) p_{i 0} \Delta t+o(\Delta t),  \tag{13}\\
r_{j i}+r_{i} \Delta t \quad \text { with probability } \mu_{j}\left(k_{j}(l)\right) u\left(k_{j}(l)\right) p_{j i} \Delta t+o(\Delta t), \\
-R_{i j}+r_{i} \Delta t \quad \text { with probability } \mu_{i}\left(k_{i}(l)\right) u\left(k_{i}(l)\right) p_{i j} \Delta t+o(\Delta t), \\
r_{i} \Delta t \quad \text { with probability } 1-\left(\lambda p_{0 i}+\mu_{i}\left(k_{i}(l)\right) u\left(k_{i}(l)\right)\right) \Delta t+o(\Delta t), \\
j=\overline{1, n}, j \neq i
\end{array}\right.
$$

where $k_{i}(l)$ is number of messages in the $i$-th QS at the $l$-th time interval, $u(x)=\left\{\begin{array}{l}1, x>0, \\ 0, x \leq 0\end{array}\right.$ - Heavyside function.

Let us find expression for expected income of system $S_{i}$ at the moment $t$. When realization of process $k(t)$ is fixed, according to (13) and denotation (5), we can write:

$$
\begin{align*}
& M\left\{d_{i l}(\Delta t) / k(t)\right\}=\left[\lambda a_{0 i} p_{0 i}+\sum_{j=1}^{n} a_{j i} p_{j i} \mu_{j}\left(k_{j}(l)\right) u\left(k_{j}(l)\right)-\right. \\
& \left.\quad-\mu_{i}\left(k_{i}(l)\right) u\left(k_{i}(l)\right) \sum_{j=0}^{n} b_{i j} p_{i j}\right] \Delta t+c_{i} \Delta t+o(\Delta t), i=\overline{1, n} \tag{14}
\end{align*}
$$

Then taking into account $m \Delta t=t$ and (3) we obtain

$$
\begin{gather*}
M\left\{V_{i}(t) / k(t)\right\}=\sum_{l=1}^{m} M\left\{d_{i l}(\Delta t) / k(t)\right\}=\left(\lambda a_{0 i} p_{0 i}+c_{i}\right) t+ \\
+\sum_{j=1}^{n} a_{j i} p_{j i} \sum_{l=1}^{m} \mu_{j}\left(k_{j}(l)\right) u\left(k_{j}(l)\right) \Delta t- \\
-\sum_{j=0}^{n} b_{i j} p_{i j} \sum_{l=1}^{m} \mu_{i}\left(k_{i}(l)\right) u\left(k_{i}(l)\right) \Delta t+o(\Delta t), i=\overline{1, n} \tag{15}
\end{gather*}
$$

When $m \rightarrow \infty, \Delta t \rightarrow 0$

$$
\sum_{l=1}^{m} \mu_{j}\left(k_{j}(l)\right) u\left(k_{j}(l)\right) \Delta t \xrightarrow[\Delta t \rightarrow 0]{ } \int_{0}^{t} \mu_{j}\left(k_{j}(s)\right) u\left(k_{j}(s)\right) d s, j=\overline{1, n}(16)
$$

so

$$
\begin{align*}
M\left\{V_{i}(t) / k(t)\right\}= & v_{i 0}+\left(c_{i}+\lambda a_{0 i} p_{0 i}\right) t+\sum_{j=1}^{n} a_{j i} p_{j i} \int_{0}^{t} \mu_{j}\left(k_{j}(s)\right) u\left(k_{j}(s)\right) d s- \\
& -\int_{0}^{t} \mu_{i}\left(k_{i}(s)\right) u\left(k_{i}(s)\right) d s \sum_{j=0}^{n} b_{i j} p_{i j}, i=\overline{1, n} \tag{17}
\end{align*}
$$

Making average by $k(t)$ and taking into account condition of normalization $\sum_{k} P(k(t)=k)=1$ for expected income of system $S_{i}$ we will have

$$
\begin{gather*}
v_{i}(t)=M\left\{V_{i}(t)\right\}=v_{i 0}+\sum_{k} P(k(t)=k) M\left\{V_{i}(t) / k(t)\right\}= \\
=v_{i 0}+\left(c_{i}+\lambda a_{0 i} p_{0 i}\right) t+\sum_{k} P(k(t)=k)\left[\sum_{j=1}^{n} a_{j i} p_{j i} \int_{0}^{t} \mu_{j}\left(k_{j}(s)\right) u\left(k_{j}(s)\right) d s-\right. \\
\left.-\int_{0}^{t} \mu_{i}\left(k_{i}(s)\right) u\left(k_{i}(s)\right) d s \sum_{j=0}^{n} b_{i j} p_{i j}\right], i=\overline{1, n} \tag{18}
\end{gather*}
$$

Let system $S_{i}$ consists of $m_{i}$ identical service lines, the service rate of messages in each line equals $\mu_{i}, i=\overline{1, n}$. In this case

$$
\mu_{i}\left(k_{i}(s)\right) u\left(k_{i}(s)\right)=\left\{\begin{array}{c}
\mu_{i} k_{i}(s), k_{i}(s) \leq m_{i},  \tag{19}\\
\mu_{i} m_{i}, k_{i}(s)>m_{i},
\end{array}=\mu_{i} \min \left(k_{i}(s), m_{i}\right), \quad i=\overline{1, n}\right.
$$

Then from (18) it follows

$$
\begin{gather*}
v_{i}(t)=M\left\{V_{i}(t)\right\}=v_{i 0}+\left(c_{i}+\lambda a_{0 i} p_{0 i}\right) t+\sum_{k} P(k(t)=k) \times \\
\times\left[\sum_{j=1}^{n} \mu_{j} a_{j i} p_{j i} \int_{0}^{t} \min \left(k_{j}(s), m_{j}\right) d s-\mu_{i} \int_{0}^{t} \min \left(k_{i}(s), m_{i}\right) d s \sum_{j=0}^{n} b_{i j} p_{i j}\right], i=\overline{1, n}(20)
\end{gather*}
$$

It is natural to suppose that average of expression $\min \left(k_{j}(s), m_{j}\right)$ gives $\min \left(N_{j}(s), m_{j}\right)$, i.e.

$$
\begin{equation*}
M \min \left(k_{j}(s), m_{j}\right)=\min \left(N_{j}(s), m_{j}\right), i=\overline{1, n} \tag{21}
\end{equation*}
$$

where $N_{j}(s)=M\left\{k_{j}(s)\right\}$ is average number of messages (waiting and servicing) in system $S_{i}$ at time interval $\left[t_{0}, t_{0}+s\right], i=\overline{1, n}$. Therefore we receive from (20) the following relation

$$
\begin{align*}
v_{i}(t)=v_{i 0}+ & \left(c_{i}+\lambda a_{0 i} p_{0 i}\right) t+\sum_{j=1}^{n} \mu_{j} a_{j i} p_{j i} \int_{0}^{t} \min \left(N_{j}(s), m_{j}\right) d s- \\
& -\mu_{i} \int_{0}^{t} \min \left(N_{i}(s), m_{i}\right) d s \sum_{j=0}^{n} b_{i j} p_{i j}, i=\overline{1, n} \tag{22}
\end{align*}
$$

As

$$
\sum_{i=1}^{n} \sum_{j=1}^{n} \mu_{j} a_{j i} p_{j i} \int_{0}^{t} \min \left(N_{j}(s), m_{j}\right) d s=\sum_{i=1}^{n} \mu_{i} \int_{0}^{t} \min \left(N_{i}(s), m_{i}\right) d s \sum_{j=1}^{n} b_{i j} p_{i j}
$$

so expected income of whole network equals

$$
M\{W(t)\}=\sum_{i=1}^{n}\left[v_{i 0}+\left(c_{i}+\lambda a_{0 i} p_{0 i}\right) t-\mu_{i} b_{i 0} p_{i 0} \int_{0}^{t} \min \left(N_{i}(s), m_{i}\right) d s\right]
$$

For finding of $N_{i}(t), i=\overline{1, n}$, it can be applied recurrence by the time moments method of analysis of average values for open QN [10].

## 4. Variations of incomes in Jackson HM-network

Taking into account (3), expression (22) can be rewritten as

$$
\begin{gather*}
v_{i}(t)=v_{i 0}+\left(c_{i}+\lambda a_{0 i} p_{0 i}\right) t-\mu_{i} b_{i 0} p_{i 0} \int_{0}^{t} \min \left(N_{i}(s), m_{i}\right) d s+ \\
+\sum_{j=1}^{n}\left[\mu_{j} a_{j i} p_{j i} \int_{0}^{t} \min \left(N_{j}(s), m_{j}\right) d s-\mu_{i} a_{i j} p_{i j} \int_{0}^{t} \min \left(N_{i}(s), m_{i}\right) d s\right], i=\overline{1, n} \tag{23}
\end{gather*}
$$

From (13), (19), (9) it follows (similar as we found (11)):

$$
\begin{gather*}
M\left\{d_{i l}^{2}(\Delta t) / k(t)\right\}=\left\{\lambda a_{20 i} p_{0 i}+\mu_{i} b_{2 i 0} p_{i 0} \min \left(k_{i}(l), m_{i}\right)+\right. \\
+\sum_{\substack{j=1 \\
j \neq i}}^{n}\left[\mu_{j} a_{2 j i} p_{j i} \min \left(k_{j}(l), m_{j}\right)+\mu_{i} b_{2 i j} p_{i j} \min \left(k_{i}(l), m_{i}\right)\right] \Delta t+o(\Delta t), i=\overline{1, n} \tag{24}
\end{gather*}
$$

Besides taking into account independence of $d_{i l}(\Delta t), d_{i q}(\Delta t), l \neq q$, and (13) we can write

$$
\begin{equation*}
M\left(d_{i l}(\Delta t) d_{i q}(\Delta t) / k(t)\right)=o(\Delta t)^{2} \tag{25}
\end{equation*}
$$

Then making limit transition when $\Delta t \rightarrow 0$, from (24), (25) and consider that $m \Delta t=t$, we have

$$
\begin{aligned}
& M\left\{\left(V_{i}(t)-v_{i 0}\right)^{2} \mid k(t)\right\}=\sum_{l=1}^{m} M\left\{d_{i l}^{2}(\Delta t) / k(t)\right\}+\sum_{l=1}^{m} \sum_{\substack{q=1 \\
q \neq l}}^{m} M\left\{d_{i l}(\Delta t) d_{i q}(\Delta t) / k(t)\right\}= \\
& =\lambda a_{20 i} p_{0 i} t+\mu_{i} b_{2 i 0} p_{i 0} \int_{0}^{t} \min \left(k_{i}(s), m_{i}\right) d s+ \\
& \quad+\sum_{\substack{j=1 \\
j \neq i}}^{n}\left[\mu_{j} a_{2 j i} p_{j i} \int_{0}^{t} \min \left(k_{j}(s), m_{j}\right) d s+\mu_{i} b_{2 i j} p_{i j} \int_{0}^{t} \min \left(k_{i}(s), m_{i}\right) d s\right], i=\overline{1, n}
\end{aligned}
$$

Making average of this expression by $k(t)$ we will have

$$
\begin{gather*}
M\left\{\left(V_{i}(t)-v_{i 0}\right)^{2}\right\}=\lambda a_{20 i} p_{0 i} t+\mu_{i} b_{2 i 0} p_{i 0} \int_{0}^{t} \min \left(N_{i}(s), m_{i}\right) d s+ \\
+\sum_{\substack{j=1 \\
j \neq i}}^{n}\left[\mu_{j} a_{2 j i} p_{j i} \int_{0}^{t} \min \left(N_{j}(s), m_{j}\right)+\mu_{i} b_{2 i j} p_{i j} \int_{0}^{t} \min \left(N_{i}(s), m_{i}\right)\right], i=\overline{1, n} \tag{26}
\end{gather*}
$$

Let us find expression for $M^{2}\left\{\left(V_{i}(t)-v_{i 0}\right)\right\}$ using (23):

$$
M^{2}\left\{\left(V_{i}(t)-v_{i 0}\right)\right\}=\left\{\left(c_{i}+\lambda a_{0 i} p_{0 i}\right) t-\mu_{i} b_{i 0} p_{i 0} \int_{0}^{t} \min \left(N_{i}(s), m_{i}\right) d s+\right.
$$

$$
\begin{equation*}
\left.+\sum_{j=1}^{n}\left[\mu_{j} a_{j i} p_{j i} \int_{0}^{t} \min \left(N_{j}(s), m_{j}\right) d s-\mu_{i} a_{i j} p_{i j} \int_{0}^{t} \min \left(N_{i}(s), m_{i}\right) d s\right]\right\}^{2}, i=\overline{1, n} \tag{27}
\end{equation*}
$$

Expression for variation of income of the $i$-th network's QS can be easily found using formula $D V_{i}(t)=M\left\{\left(V_{i}(t)-v_{i 0}\right)^{2}\right\}-M^{2}\left\{\left(V_{i}(t)-v_{i 0}\right)\right\}$ and relations (26), (27).

## 5. Some optimization problems for network incomes

Let us to consider open QN with incomes that was described in the previous sections. Let us formulate two optimization problems that are associated with income's maximization of the whole network or central QS for example

$$
\begin{gather*}
\left\{\begin{array}{c}
W_{1}(T, m)=W\left(T, m_{1}, \ldots, m_{n}\right)=\frac{1}{T-t_{0}} \int_{t_{0}}^{T} \sum_{i=1}^{n}\left(M V_{i}(t)-d_{i} N_{i}(t)-E_{i} m_{i}\right) d t \rightarrow \max _{m_{1}, m_{2}, \ldots, m_{n}} \\
m_{i} \leq a_{i}, i=\overline{1, n}
\end{array}\right.  \tag{28}\\
\left\{\begin{aligned}
W_{2}(T, m)=W_{2}\left(T, m_{1}, \ldots, m_{n}\right)= & \frac{1}{T-t_{0}} \int_{t_{0}}^{T}\left(v_{n}(t)-d_{n} N_{n}(t)-E_{n} m_{n}\right) d t \rightarrow \max _{m_{1}, m_{2}, \ldots, m_{n}} \\
& m_{i} \leq a_{i}, i=\overline{1, n}
\end{aligned}\right. \tag{29}
\end{gather*}
$$

where $a_{i}$ - set numbers that present range in which number of service lines in the $i$ th QS can change; $E_{i}$ - coasts for keeping of one service line; $d_{i}$ - coasts for keeping of one message in the $i$-th system, $i=1, n$. From relation (20) it is evident that (28) and (29) can be reduced to the next problems

$$
\begin{gather*}
\left\{\frac{1}{T-t_{0}} \int_{t_{0}}^{T} \sum_{i=1}^{n}\left(\mu_{i} b_{i 0} p_{i 0} \int_{t_{0}}^{t} \min \left(N_{i}(s), m_{i}\right) d s+d_{i} N_{i}(t)+E_{i} m_{i}\right) d t \rightarrow \min _{m_{1}, m_{2}, \ldots, m_{,}}\right.  \tag{30}\\
m_{i} \leq a_{i}, i=\overline{1, n}
\end{gather*}\left\{\begin{array}{c}
\frac{1}{T-t_{0}} \int_{t_{0}}^{T}\left[\sum_{j=1}^{n} \mu_{j} a_{j n} p_{j n}^{t} \int_{t_{0}}^{t} \min \left(N_{j}(s), m_{j}\right) d s-\mu_{n} \int_{t_{0}}^{t} \min \left(N_{n}(s), m_{n}\right) d s \sum_{j=0}^{n} b_{n j} p_{n j}-\right.  \tag{31}\\
\left.-\left(d_{n} N_{n}(t)+E_{n} m_{n}\right)\right] d t \rightarrow \max _{m_{1}, m_{2}, \ldots, m_{n}} \\
m_{i} \leq a_{i}, i=\overline{1, n}
\end{array}\right.
$$



Fig. 1. Topology of HM-network

Integrals in problems (30) and (31) can be found by applying method of trapeziums two times sequentially.

## 6. Modelling example

Let us consider open HM-network with central system, (Fig. 1). Network is discribed with the next parametrs: count of QS in network $n=20, \lambda=15$; the starting vector of service line's count in network systems consists of elements $m_{i}=1$, $N_{i}(0)=5, \quad a_{i}=10, i=\overline{1,20}$; probobilities of message transitions between network's systems are equal $p_{0 i}=0.063 q^{i-1}, q=0.975, i=\overline{1,20}$, (they are found with a help of formula of geometric series and condition $\sum_{i=1}^{20} p_{0 i}=1$, $\quad p_{i 0}=p_{i 20}=0.5, i=\overline{1,19}$, $p_{200}=0.5, \quad p_{20 i}=0.026, i=\overline{1,19}$, other $p_{i j}=0, i, j=\overline{1,19}$. In this case vector of message arrival rates in each system consists of components $\lambda p_{0 i}$ and equals (1.219, 1.195, 1.172, 1.149, 1.128, 1.106, 1.086, 1.065, 1.045, 1.026, 1.01, 0.989, $0.971,0.954,0.937,0.92,0.904,0.888,0.873,10.401)$. Service times of messages in each system's line are distributed with exponential law with parametrs $\mu_{i}$ that are equal ( $2.119,2.695,2.07,1.349,1.728,2.606,1.386,1.165,2.045,2.426,1.707$, $1.989,2.371,1.754,2.037,2.32,1.204,1.588,0.973,11.8)$; vector with components $d_{i}$ looks like $(0.25,0.28,0.26,0.5,0.8,0.3,0.17,0.5,0.10,0.7,0.10,0.11$, $0.17,0.18,0.10,0.80,0.2,0.51,0.4,0.14)$ and vector with components $E_{i}-(0.10$, $0.19,0.17,0.8,0.13,0.3,0.5,0.2,0.11,0.4,0.11,0.16,0.7,0.18,0.13,0.15,0.15$ $, 0.20,0.12,0.10)$. Let also e.v. are equal $a_{j 20}=27, j=1,18,19,20$, $a_{j 20}=31, j=2,3,4,5,9,10,11, \quad a_{j 20}=19, \quad j=8,15,16,17$, $a_{j 20}=21, j=6,7,12,13,14 ; \quad b_{j 0}=11, j=11,20, \quad b_{j 0}=13, \quad j=5,7,9,16$,
$b_{j 0}=15, j=2,6,13,15,18, \quad b_{j 0}=17, j=3,4,14,19, \quad b_{j 0}=19, \quad j=1,8,10,12,17$;
$b_{n j}=12, j=3,14,15,19, \quad b_{n j}=14, j=5,20, \quad b_{n j}=16, \quad j=2,4,9,10,18$,
$b_{n j}=18, j=7,13,17, b_{n j}=11, j=1,6,8,1,12,16 ; \quad t_{0}=0, T=50$.
Optimization problems (30) and (31) were solved by method of full search. Solution of problem (30) is $m_{i}^{*}=1, i=\overline{1,19}, m_{20}^{*}=2$ and value of optimization criteria in this case is $W_{1}\left(50, m^{*}\right)=9.138$. And solution of optimization problem (31) is $m_{i}^{*}=1, i=\overline{1,19}, i \neq 4, \quad m_{4}^{*}=2, \quad m_{20}^{*}=3$, and value of optimization criteria in this case is $W_{2}\left(50, m_{i}^{*}\right)=-4.94$.

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